

# Cray Reveal Webinar: A Tool to Help Porting to Manycore

Heidi Poxon
Technical Lead & Sr. Manager, Performance Tools
Cray Inc.

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#### **Future Architecture Directions**



# Nodes are becoming more parallel

- More processors per node
- More threads per processor
- Vector lengths are getting longer
- Memory hierarchy is becoming more complex
- Scalar performance is not increasing and will start decreasing

#### For the next decade, HPC systems will have the same basic architecture:

- Message passing between nodes
- Multithreading within the node (pure MPI will not do)
- Vectorization at the lowest level (SSE, AVX, GPU, Phi)

ANALYZE

# **Future Application Directions**



- Threading on node as well as vectorization is becoming more important – need more parallelism exploited in applications due to increasing number of cores and threads
- Current petascale applications are not structured to take advantage of these architectures
  - Currently 80-90% of applications use a single level of parallelism
    - MPI or PGAS between cores of the MPP system
  - Looking forward, application developers are faced with a significant task in preparing their applications for the future
    - Codes must be converted to use multiple levels of parallelism
    - More complex memory hierarchies will require user intervention to achieve good performance

# Three Levels of Parallelism Required

- Developers will continue to use MPI between nodes or sockets
- 2. Developers must address using a shared memory programming paradigm on the node
- 3. Developers must vectorize low level looping structures

While there is a potential acceptance of new languages for addressing all levels directly. Most developers cannot afford this approach until they are assured that the new language will be accepted and the generated code is within a reasonable performance range

# When to Move to a Hybrid Programming Model



#### When code is network bound

- Look at collective time, excluding sync time: this goes up as network becomes a problem
- Look at point-to-point wait times: if these go up, network may be a problem

#### When MPI starts leveling off

- Too much memory used, even if on-node shared communication is available
- As the number of MPI ranks increases, more off-node communication can result, creating a network injection issue
- When contention of shared resources increases
- When you want to exploit heterogeneous nodes

# **Approach to Adding Parallelism**



#### 1. Identify key high-level loops

- Determine where to add additional levels of parallelism
  - Assumes MPI application is functioning correctly on X86
  - Find top serial work-intensive loops (perftools + CCE loop work estimates)

# 2. Perform parallel analysis, scoping and vectorization

- Split loop work among threads
  - Do parallel analysis and restructuring on targeted high level loops
  - Use Reveal + CCE for scoping, loopmark and source browsing

# 3. Add OpenMP layer of parallelism

- Insert OpenMP directives (with Reveal directive building assistance)
  - Run on X86 to verify application and check for performance improvements
- 4. Analyze performance for further optimizations, specifically vectorization of innermost loops

# Challenges



- Investigate parallelizability of high level looping structures
  - Often times one level of loop is not enough, must have several parallel loops
    - Need a large number of loop iterations to feed the GPU threads
  - User must understand which high level DO loops have independent iterations
  - Without tools, variable scoping of high level loops is very difficult
    - Loops must be more than independent, their variable usage must adhere to private data local to a thread or global shared across all the threads
    - Independence can be complicated to understand (and even runtime dependent)
- Investigate vectorizability of lower level DO loops

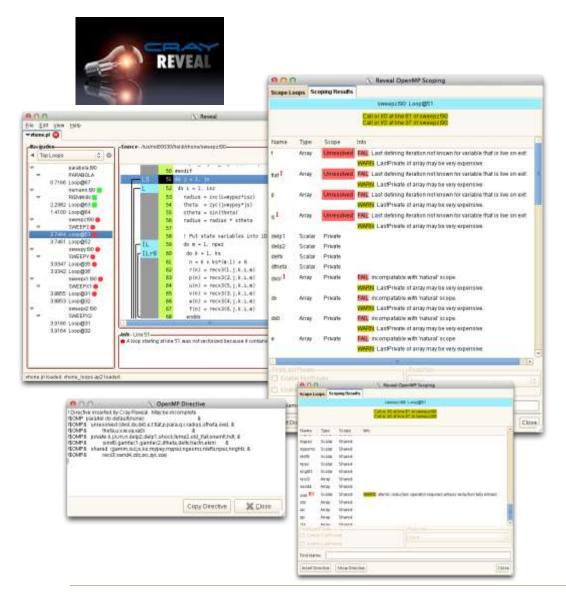
# The Problem – How Do I Parallelize This Loop?

- How do I know this is a good loop to parallelize?
- What prevents me from parallelizing this loop?
- Can I get help building a directive?

```
subroutine sweepz
do j = 1, js
do i = 1, isz
  radius = zxc(i+mypez*isz)
  theta = zyc(j+mypey*js)
  do m = 1, npez
   do k = 1, ks
   n = k + ks*(m-1) + 6
    r(n) = recv3(1,j,k,i,m)
    p(n) = recv3(2,j,k,i,m)
    u(n) = recv3(5,j,k,i,m)
    v(n) = recv3(3,j,k,i,m)
    w(n) = recv3(4,j,k,i,m)
    f(n) = recv3(6,j,k,i,m)
   enddo
  enddo
  call ppmlr
  do k = 1, kmax
    n = k + 6
    xa(n) = zza(k)
    dx(n) = zdz(k)
    xa0(n) = zza(k)
    dx0(n) = zdz(k)
    e(n) = p(n)/(r(n)*qamm)+0.5 &
        (u(n)**2+v(n)**2+w(n)**2)
  enddo
  call ppmlr
enddo
enddo
```

```
subroutine ppmlr
call boundary
call flatten
call paraset(nmin-4, nmax+5, para, dx, xa)
call parabola(nmin-4,nmax+4,para,p,dp,p6,p1,flat)
call parabola(nmin-4,nmax+4, para,r,dr,r6,r1,flat)
call parabola(nmin-4,nmax+4,para,u,du,u6,u1,flat)
call states(pl,ul,rl,p6,u6,r6,dp,du,dr,plft,ulft,&
           rlft,prqh,urqh,rrqh)
call riemann(nmin-3,nmax+4,gam,prgh,urgh,rrgh,&
            plft,ulft,rlft pmid umid)
call remap ← contains more calls
call volume(nmin,nmax,ngeom,radius,xa,dx,dvol)
call remap ← contains more calls
return
End
```

# Simplifying the Task with Reveal



- Navigate to relevant loops to parallelize
- Identify parallelization and scoping issues
- Get feedback on issues down the call chain (shared reductions, etc.)
- Optionally insert parallel directives into source
- Validate scoping correctness on existing directives

COMPUTE

# **Using Reveal with Performance Statistics**



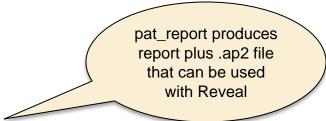
# Optionally create loop statistics using the Cray performance tools to determine which loops have the most work

- Helps identify high-level serial loops to parallelize
  - Based on runtime analysis, approximates how much work exists within a loop
- Provides the following statistics
  - Min, max and average trip counts
  - Inclusive time spent in loops
  - Number of times a loop was executed

# **Collecting Loop Work Estimates**



- Load PrgEnv-cray module (must use CCE)
- Load perftools module
- Compile AND link with –h profile\_generate
  - cc –h profile\_generate –o my\_program my\_program.c
- Instrument binary for tracing
  - pat\_build –w my\_program
- Run application
- Create report with loop statistics
  - pat\_report my\_program.xf > loops\_report



# **Example Report** – Inclusive Loop Time



```
Table 2: Loop Stats by Function (from -hprofile generate)
         Loop | Loop | Loop | Function=/.LOOP[.]
   Loop
   Incl
       | Hit
                    Trips
                          | Trips
                                         | PE=HIDE
                                  | Trips
  Time
                    Avg | Min |
                                    Max
  Total
                        25 I
                            0 | 25 | sweepy .LOOP.1.li.33
 8.995914
           100 I
                    25 | 0 | 25 | sweepy_.LOOP.2.li.34
| 8.995604 | 2500 |
                    25 | 0 | 25 | sweepz .LOOP.05.1i.49
| 8.894750 |
          50 I
                    25 | 0 | 25 | sweepz .LOOP.06.1i.50
I 8.894637 I
          1250 |
| 4.420629 |
                            0 | 25 |sweepx2 .LOOP.1.li.29
             50 I
                       25 I
| 4.420536 |
          1250
                    25 | 0 | 25 | sweepx2 .LOOP.2.1i.30
          50 |
                            0 |
| 4.387534 |
                       25 |
                                     25 |sweepx1 .LOOP.1.li.29
| 4.387457 |
                                     25 |sweepx1 .LOOP.2.li.30
          1250
                     25 |
                             0 1
| 2.523214 | 187500 |
                                      107 | riemann .LOOP.2.1i.63
                    107 |
                                0 |
| 1.541299 | 20062500 |
                                       12 | riemann .LOOP.3.1i.64
                    12 |
                                0 |
| 0.863656 |
           1687500 I
                                      108 |parabola .LOOP.6.li.67
                       104 |
                                0 |
```

#### **How to Use Reveal**



# Generate a program library for your application with CCE

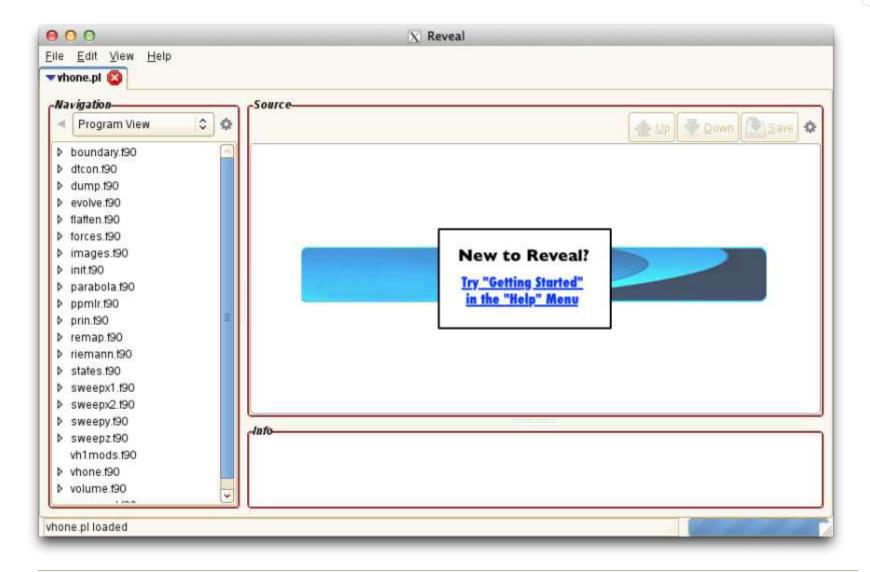
- > cc -h pl=himeno.pl -hwp himeno.c
- > ftn -h pl=vhone.pl -hwp file1.f90

Optionally add whole program analysis for more aggressive inlining

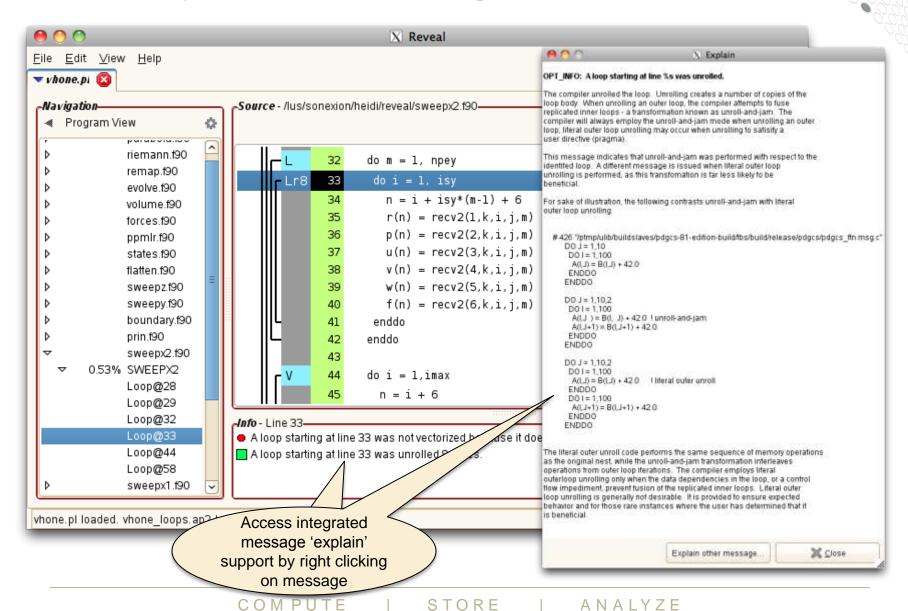
#### Launch Reveal

- > module load perftools
- Use with compiler information only (no need to run program):
  - > reveal vhone.pl
- Use with compiler + loop work estimates (include performance data)
  - > reveal vhone.pl vhone\_loops.ap2

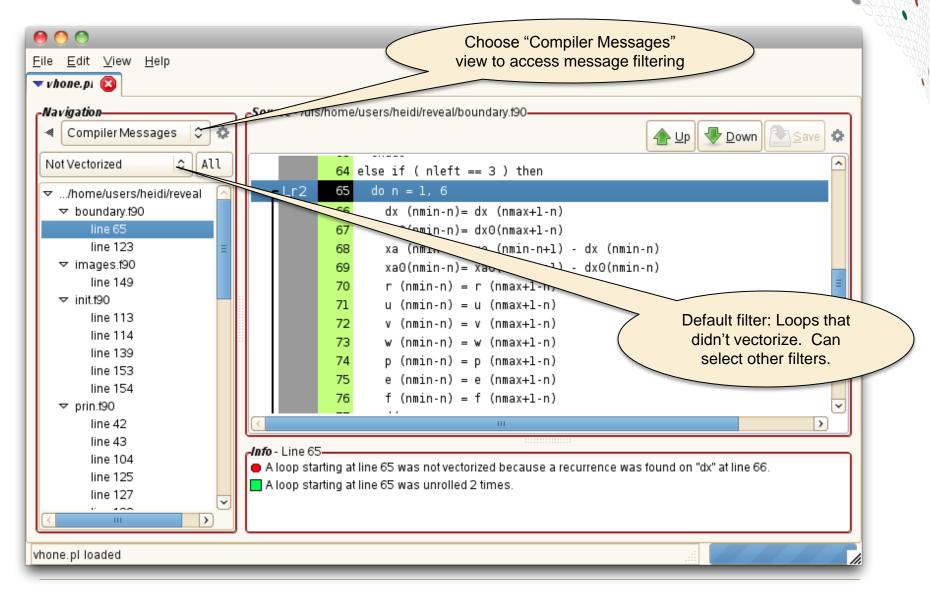
# **Browse Source and Compiler Optimizations**



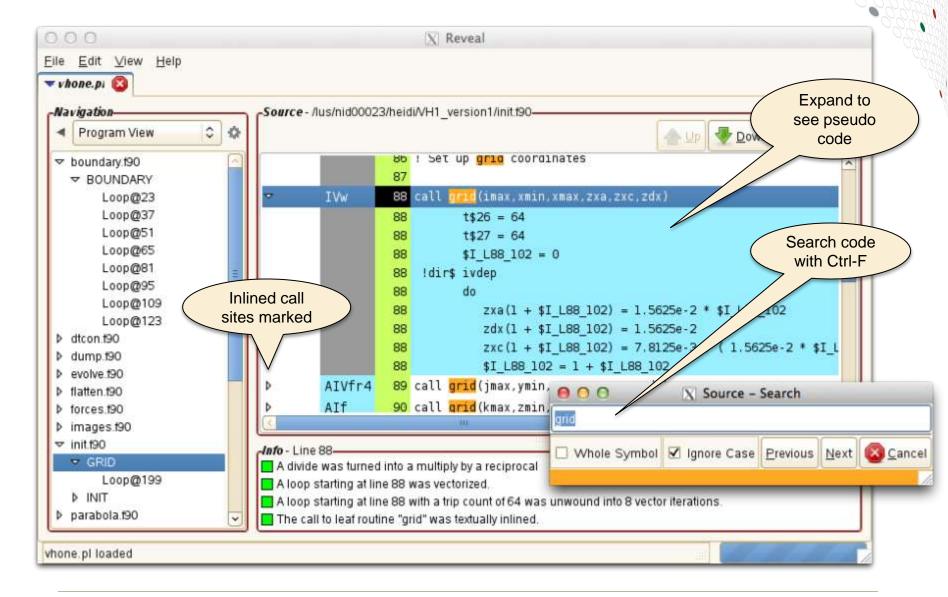
# **Access Cray Compiler Message Information**



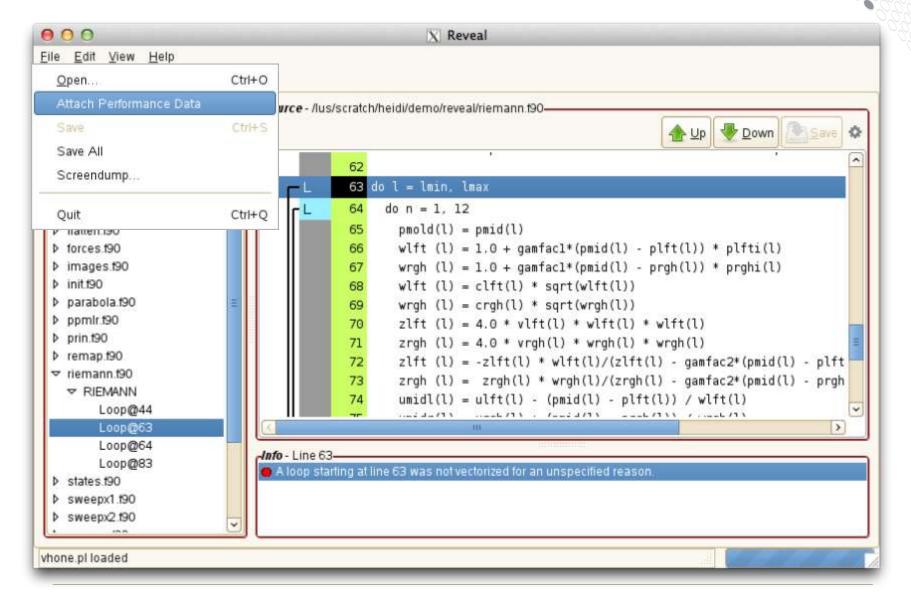
# **Navigate Code via Compiler Messages**



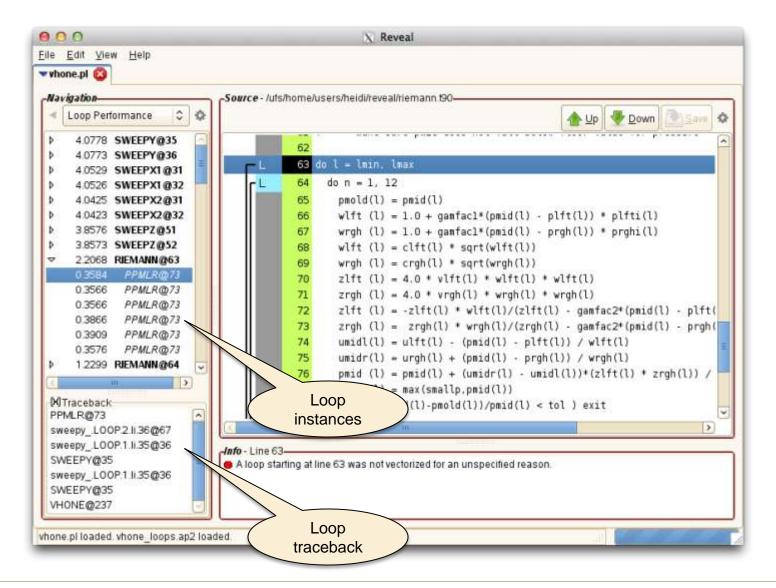
#### **View Pseudo Code for Inlined Functions**



# **Add Performance Data to Find Top Loops**

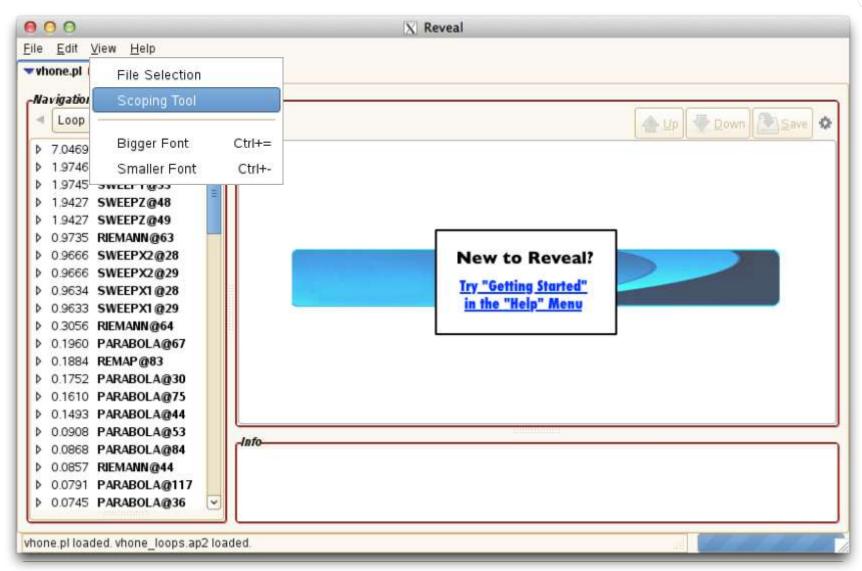


# **View Loops through Call Chain**

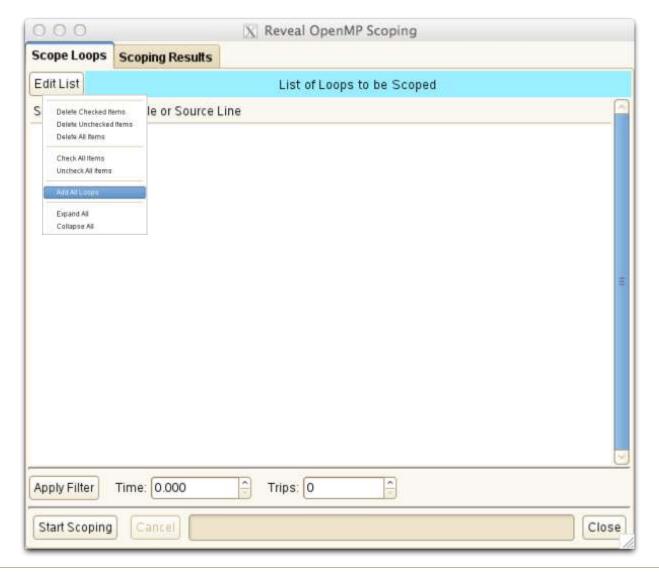


# **Scope Top Time Consuming Loops**



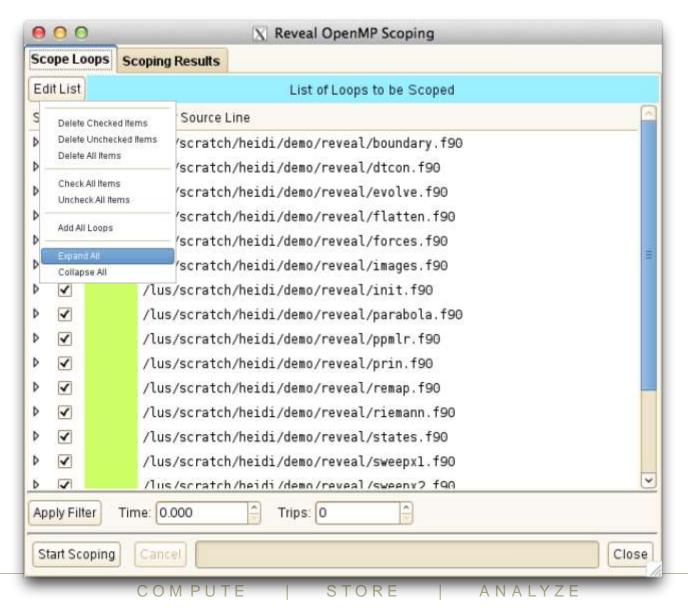


# **Include All Loops as Initial Candidates**

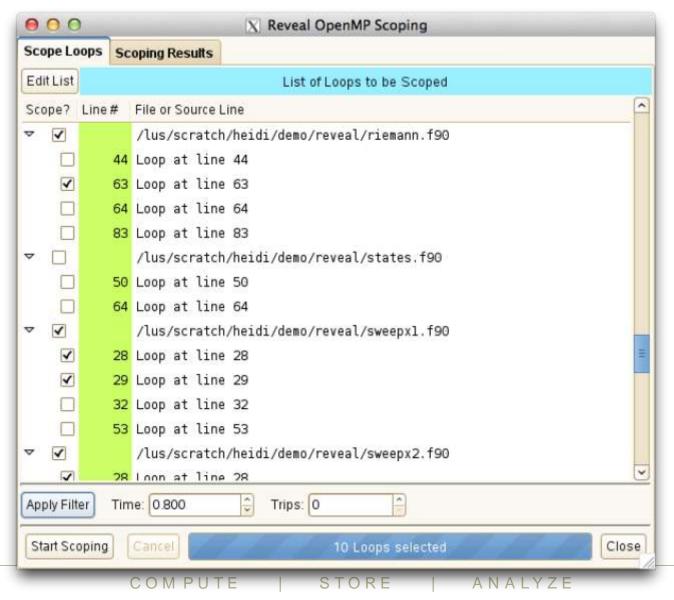




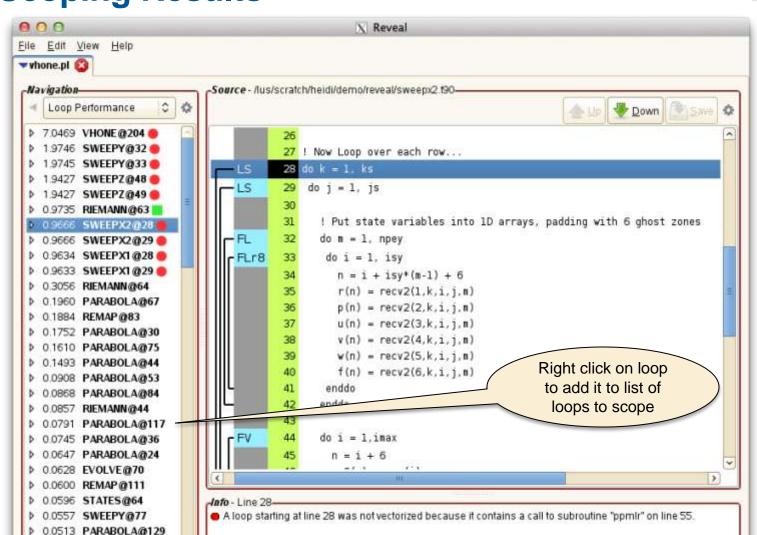
# Include All Loops as Initial Candidates (2)



# **Apply Filter to Select Only Top Loops**

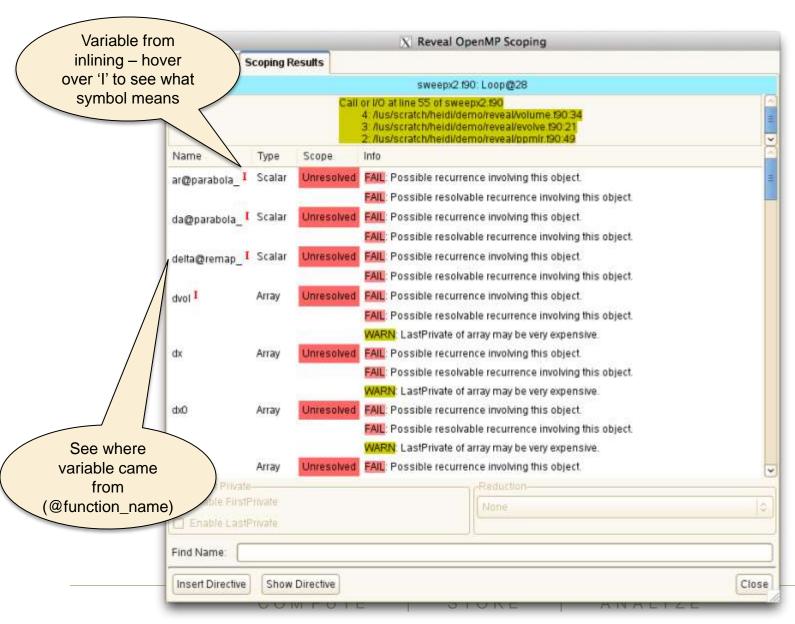


# **View Scoping Results**

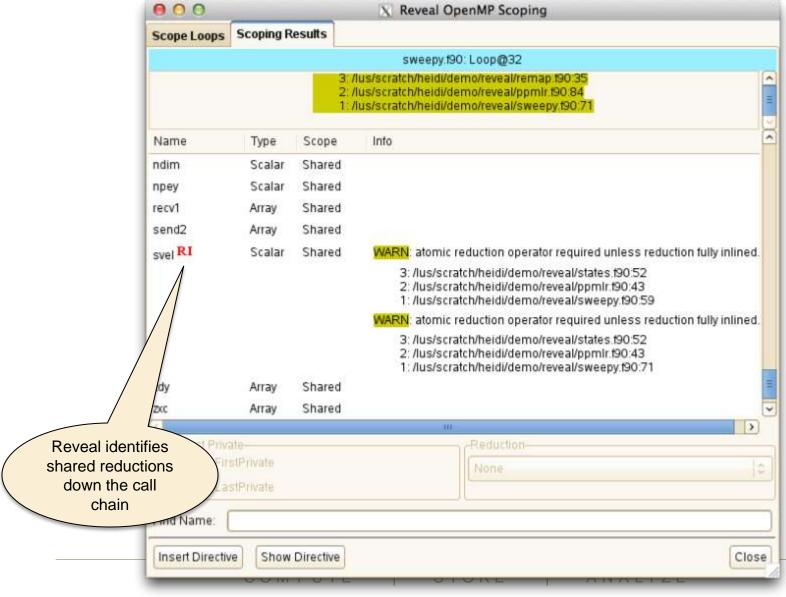


whone pl loaded whone loops ap2 loaded

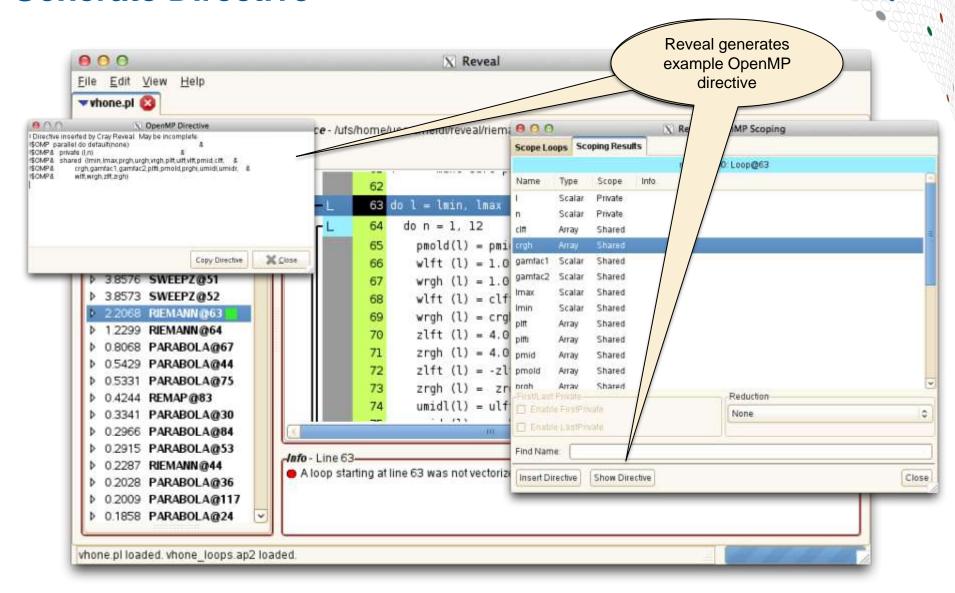
# Reveal Gives Feedback on Scoping Results



#### **Reveal Points Out Parallelization Issues**

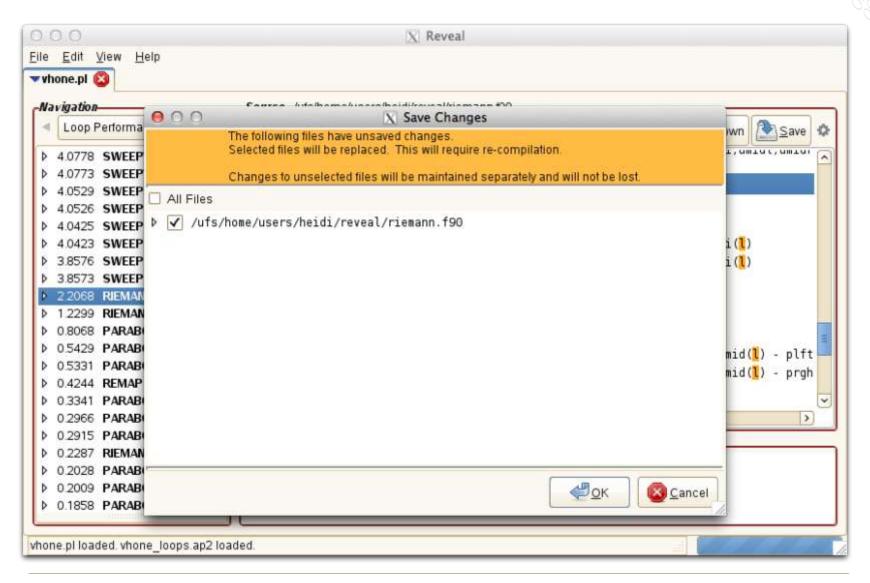


#### **Generate Directive**



# **Optionally Insert Directive Into Source**





#### **Reveal Inserts Directive Into Source**



```
! Directive inserted by Cray Reveal. May be incomplete.
!$OMP parallel do default(none)
        unresolved (dvol,dx,dx0,e,f,flat,p,para,q,r,radius,svel,u,v,w,
!$OMP&
!$OMP&
                 xa,xa0)
        !$OMP&
                 one 1, hdt, sinxf0, qamfac1, qamfac2, dtheta, deltx, fractn, &
!$OMP&
!SOMP&
                 ekin)
                            \ks, mypey, ndim, ngeomy, nlefty, npey, nrighty,
!$OMP&
        shared
                (gamm, isy,
                 recv1, send2,
!$OMP&
                                zxc,zya)
do k = 1, ks
do i = 1, isy
  radius = zxc(i+mypey*isy)
                                                 th 6 ghost zones
  ! Put state variables into 1D arrays, padd
  do m = 1, npey
   do j = 1, js
    n = j + js*(m-1) + 6
    r(n) = recv1(1,k,j,i,m)
    p(n) = recv1(2,k,j,i,m)
    u(n) = recv1(4,k,j,i,m)
    v(n) = recv1(5,k,j,i,m)
    w(n) = recv1(3,k,j,i,m)
    f(n) = recv1(6,k,j,i,m)
   enddo
  enddo
  do j = 1, jmax
    n = j + 6
```

Reveal generates
OpenMP directive
with illegal clause
marking variables
that need addressing

# Resolve Private Array Concerns for dvol, etc.

For OpenMP these need to be made task\_private

!\$omp threadprivate(dvol,dx,dx0,e,f,flat,p,para,q,r,radius,theta,stheta,u,v,w,xa,xa0)

#### **Resolve Shared Reductions**



#### Original

```
hdt = 0.5*dt
do n = nmin-4, nmax+4
  Cdtdx (n) = sqrt(gam*p(n)/r(n))/(dx(n)*radius)
  svel = max(svel,Cdtdx(n))
  Cdtdx (n) = Cdtdx(n)*hdt
  fCdtdx(n) = 1. - fourthd*Cdtdx(n)
enddo
```

#### Restructured – One Approach

```
hdt = 0.5*dt
!$omp critical
do n = nmin-4, nmax+4
  Cdtdx (n) = sqrt(gam*p(n)/r(n))/(dx(n)*radius)
  svel = max(svel,Cdtdx(n))
  Cdtdx (n) = Cdtdx(n)*hdt
  fCdtdx(n) = 1. - fourthd*Cdtdx(n)
enddo
!$omp end critical
```

For OpenMP need to have a critical region around setting of svel



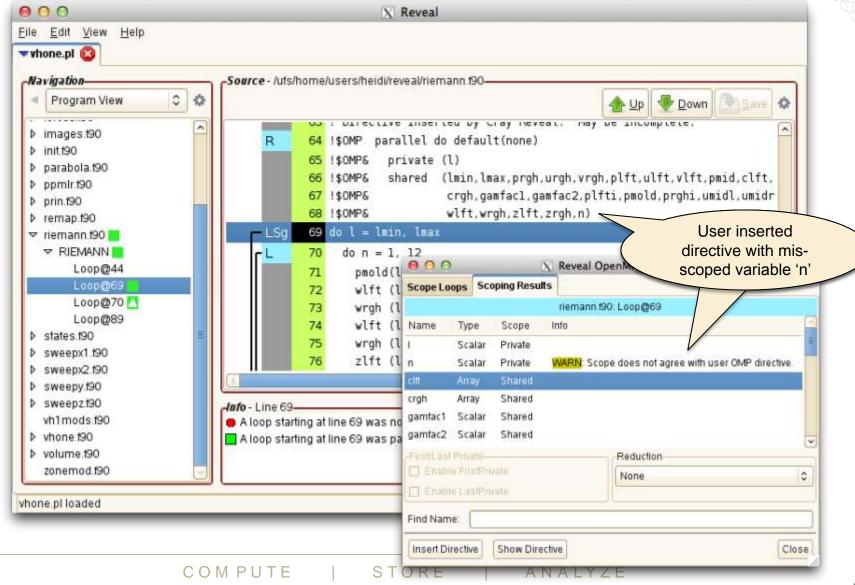


```
hdt = 0.5*dt
do n = nmin-4, nmax+4
   Cdtdx (n) = sqrt(gam*p(n)/r(n))/(dx(n)*radius)
   svel = max(svel,Cdtdx(n))
   Cdtdx (n) = Cdtdx(n)*hdt
   fCdtdx(n) = 1. - fourthd*Cdtdx(n)
enddo
```

#### Restructured – Better Approach

```
hdt = 0.5*dt
Svel0 = 0.0
do n = nmin-4, nmax+4
   Cdtdx (n) = sqrt(gam*p(n)/r(n))/(dx(n)*radius)
   svel0(n) = max(svel(n),Cdtdx(n))
   Cdtdx (n) = Cdtdx(n)*hdt
   fCdtdx(n) = 1. - fourthd*Cdtdx(n)
Enddo
!$omp critical
Do n = nmin-4, nmax +4
   Svel = max(svel0(n),svel)
Enddo
!$omp end critical
```

#### **Use Reveal to Validate User Inserted Directives**



# VH1 – Astrophysics Code



- VH1 is written with high level loops and complex decision processes.
- Ported to hybrid MPI + OpenMP using Reveal
- Reveal was able to identify
  - storage conflicts
  - private variables in modules
  - reductions down the call chain that require critical regions
- Scoping was performed in seconds where it would have taken weeks to get correct without Reveal

### S3D - Structured Cartesian Mesh Flow Solver



- S3D, a pure MPI program, was converted to a hybrid multicore application suited for a multi-core node with or without an accelerator.
- When the work was started, Reveal did not exist.
- Once Reveal was available, it was instrumental in identifying bugs in the scoping of extremely large loops (3000 lines of Fortran).
- There are both OpenMP and OpenACC versions of S3D that run well on both OpenMP systems and on the Titan Cray XK7 machine at Oak Ridge National Laboratory.

# **Summary**



- Reveal can be used to simplify the task of adding OpenMP to MPI programs
- Can be used as a stepping stone for codes targeted for nodes with higher core counts and as the first step in adding OpenACC to applications to for execution on GPUs

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